

The Ultimate Baseball Road Trip and Algorithm

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Combinatorial Optimization

- Optimize a combination from a finite set of objects
- Define a set of constraints that a solution must satisfy
- Define criteria to compare one solution to another

Major League Baseball (MLB) schedule

- 30 teams each play 81 home games in a season, between April and September
- $30 \times 81 = 2430 \text{ total games}$
- Find the most efficient driving route to attend one game at each stadium
- Primary optimization: Minimize total number of days (trip duration)
- Secondary: Minimize driving time (*travel* duration)

Constraints

Our problem space has 2,494 decision variables and approximately 120,000 constraints

https://harvardsportsanalysis.files.wordpress.com/2011/06/shortest_possible_baseball_road_trip1.pdf

(1)
$$B \leq S_{n,b} + (1 - X_{n,b}) * M \quad \forall b \in B, n \in N$$

(2)
$$E \ge S_{n,b} + (1 - X_{n,b}) * M \quad \forall b \in B, n \in N$$

(3)
$$\sum_{n \in N} X_{n,b} = 1 \quad \forall b \in B$$

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$$(4) |X_{n,b} + X_{n',b'} \le 1 |\forall b, b' \in B, n, n' \in N | such that |S_{n,b} - S_{n',b'}| \le T_{b,b'} + D$$

(5)
$$\sum_{n \in N} X_{n,b} * S_{n,b} \le (1 - F_b) * M + \sum_{n' \in N} X_{n',b'} * S_{n',b'} \quad \forall b, b' \in B$$

(6)
$$(1-L_b)*M+\sum_{n\in\mathbb{N}}X_{n,b}*S_{n,b}\leq\sum_{n'\in\mathbb{N}}X_{n',b'}*S_{n',b'} \quad \forall b,b'\in B$$

(7)
$$\sum_{b \in B} F_b = 1$$

(8)
$$\sum_{b \in B} L_b = 1$$

(9)
$$P \ge T_{b,b'} - (2 - F_b - L_{b'}) * M \quad \forall b, b' \in B$$

Domain-Specific Heuristics

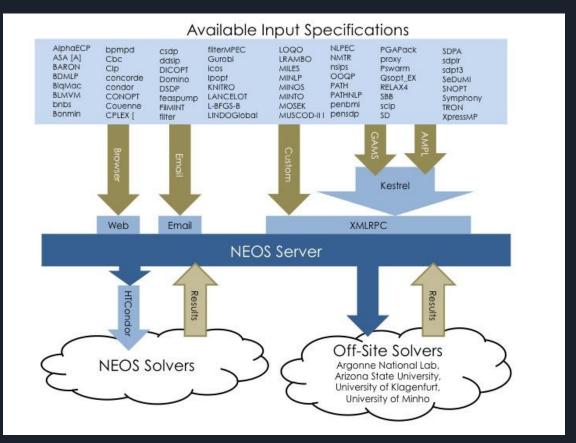
- Isolated West Coast stadiums
- Reasonable upper bound
- All-Star Break (early July)

Creating Subsets of the Problem Space

- 1. Identify a date range where the seven West Coast stadiums can be visited in no more than nine days
- 2. Select all games from other stadiums that are scheduled within 26 days of either end of the West Coast date range
- 3. Remove any games which are on a different side of the All-Star Break from the West Coast date range
- 4. Repeat steps 1-3 with a different West Coast date range

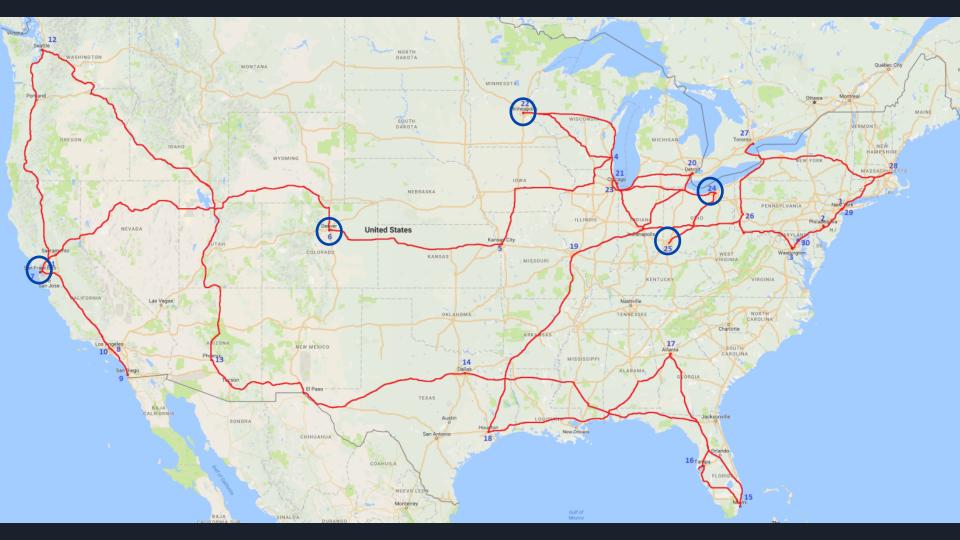
The larger subsets will have about 700 decision variables and 15,000 constraints (~20x reduction in computational complexity)

NEOS (Network-Enabled Optimization System)



```
GUB cover cuts applied: 8
Clique cuts applied: 73
Cover cuts applied: 28
Implied bound cuts applied: 39
Flow cuts applied: 27
Mixed integer rounding cuts applied: 151
Zero-half cuts applied: 40
Lift and project cuts applied: 61
Gomory fractional cuts applied: 5
Root node processing (before b&c):
 Real time
                = 12.75 sec. (2154.94 ticks)
Parallel b&c, 4 threads:
 Real time
                 = 178.28 sec. (36540.97 ticks)
 Sync time (average) = 27.16 sec.
 Wait time (average) = 0.01 sec.
Total (root+branch&cut) = 191.03 sec. (38695.90 ticks)
Solution pool: 16 solutions saved.
MIP - Integer optimal, tolerance (0.0001/1e-06): Objective = 4.9269000000e+04
Current MIP best bound = 4.9266243190e+04 (gap = 2.75681, 0.01%)
Solution time = 191.03 sec. Iterations = 844453 Nodes = 11742 (1)
Deterministic time = 38695.91 ticks (202.56 ticks/sec)
```

	CPLEX> Incumbent solution		
V	/ariable Name S	olution Value	
e	endTime	340715.000000	
s	startTime	291635.000000	
t	timeToStart	189.000000	
N	NYYJul21	1.000000	
P	PHIJul21	1.000000	
W	NASJul22	1.000000	
M	MILJul23	1.000000	
K	(CRJul24	1.000000	
	OLJul25	1.000000	
9	SFGJul27	1.000000	
9	SDPJul28	1.000000	
L	LAAJul29	1.000000	
	DAKJul30	1.000000	
9	SEAAug01	1.000000	
L	LADAug03	1.000000	
A	ARIAug04	1.000000	
T	TEXAug06	1.000000	
M	1IAAug08	1.000000	
T	TBRAug09	1.000000	
A	ATLAug10	1.000000	
H	HOUAug12	1.000000	
2	STLAug13	1.000000	
	DETAug14	1.000000	
	THCAug15	1.000000	
M	4INAug16	1.000000	
	IWSAug17	1.000000	
	LEAug18	1.000000	
	INAug19	1.000000	
P	PITAug20	1.000000	
	TORAug21	1.000000	
N	NYMAug22	1.000000	
E	BOSAug23	1.000000	
E	BALAug24	1.000000	
i	isFirstNYY	1.000000	
	isLastBAL	1.000000	
	All other variables in the range 1-538 are 0.		
	CPLEX>		





N0000000000000

JULY 21 PHIL

Tonight's game between the Phillies a will be played tomorrow, July 22, as part of a (game 49 on ticket) scheduled for a 6:05 p.m. on ticket) will be moved from 1:35 p.m. to 12 and for the second game will open approxim

All fans holding tickets to tonight's ga original tickets. Fans unable to attend that ga season home game in 2018.

The PBCO Phillies Scicks will be given children 14 and under attending the 6:05 p.m day, originally scheduled for today, has been

Both games will be televised on NBC

Philadelphia Phillies @ @Phillies

Tonight's #Phillies game against the San Diego Padres has been rescheduled and will be played as part of a split-admission doubleheader tomorrow, July 22....

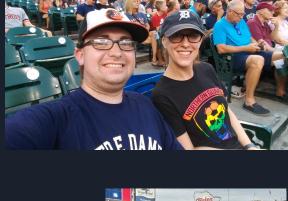


Trip Stats

- 30 MLB games in 36 days
- 16,935 miles driven
- Every major US interstate (multiples of 5)
- 34 states + DC and Ontario
- 49 gas stops
- 574 players (77% of major league rosters)
- Four MLB debuts
- Orioles record: 9-23



- 2 guest appearances on FanGraphs podcast Effectively Wild (episodes 1169 and 1263)
- @MLBRoadTrip on Twitter with ballpark photos, stadium food selections and road trip playlists
- Interviews with WCBS radio in New York, Shepherd Express in Milwaukee, Fast Company and Technical ly Baltimore
- Meetups with family, friends, and fans all over the country









Source Code

https://github.com/mountm/mlbroadtrip

Special Thanks

- Ben Blatt & Harvard Sports Analysis Collective
- neos-server.org
- Catalyte
- Effectively Wild (@EWPod)
- Interstate Highway System
- Assistant driver Dave Mountain (7/22, 8/2-8/10)



Source code at github.com/mountm/mlbroadtrip

Optimizing for Travel Distance with the Backtracking Approach

Backtracking approach

- Incrementally build a candidate solution
- Abandon your candidate (backtrack) as soon as you can confirm that the candidate is invalid (unable to be extended)
- Not a formal algorithm, more of a metaheuristic

Two Questions

- What does a partial candidate solution looks like?
- How do we know if a candidate solution is valid?

Backtracking pseudocode, where the entry point is bt(root(P))

```
procedure bt(c)

if reject(P, c) then return

if accept(P, c) then output(P, c)

s \leftarrow first(P, c)

while s \neq \Lambda do

bt(s)

s \leftarrow next(P, s)
```

 Λ is the *null candidate*

Let's define our procedural parameters!

P - an array of games in the schedule, starting from the beginning point of our time-optimized trip, ending at the final allowed day, sorted by start time

root(P) - the first game in the time-optimized solution

reject(P, c) - verify that each game in the candidate c can be reached from the previous one and that no stadium is visited twice

accept(P, c) - verify that the length of the candidate array is 30

first(P, c) - From P, select the next game after the final game in c, and add it to c (or return Λ)

next(P, s) - From P, select the next game after the final game in c and replace the final game in c

output(P, c) - Calculate the total driving distance for c. If no previous solution was found, accept this as the best solution. Otherwise, keep the solution with the shortest distance needed.

Improvements to Consider

- The reject function should err on the side of caution
- But it should be as aggressive as possible
- Effective equivalence of partial candidates for earlier pruning
- Trading memory usage for execution time (e.g. parallelization)